Encryption Using Timing Clock

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Abstract— In this research will be work to encryption algorithm based on a new idea to encrypt text data, which depends on the architectural and mathematical operations and then generate the cipher text output the these operations where we will be the removal of the letters and replace letters, numbers and symbols depending on the specific algorithm indicated the pulses time, finally this research It includes clarification of a new way to encrypt data using architectural thought to get rid the frequencies of letters, in addition to the algorithm will be the launch of its name (Clock encryption).

Index Terms— Computer Security, Data Encryption, Clock, CTE, Symmetric, Substitution, Key, Frequency, shifting, Characters, Cipher.

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I INTRODUCTION

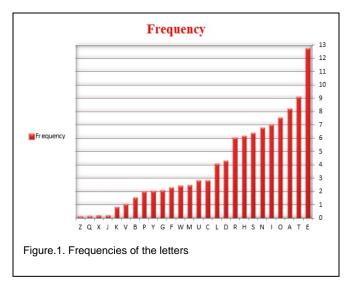
ed the increasing growth of multimedia applications on a communication networks to the increase in the need to provide efficient modalities Which works to protect data and private ownership of the individual and therefore had to be a means to act providing for the security of these media To protect it from theft hacking and tampering with counterfeiting and dissemination of sensitive information.

Hence the need for the provision of data security, and this means the encryption flag who cares to provide protection for the storage and transfer the data through the use of a secret key, therefore encryption still successful way to protect data stored that sent across the network, but with an increase in high traffic areas networks and information network internet world, it has become difficult to keep these data, particularly that it always be accessible through the Internet in the clear wording sends that Doubt The attention of the intruder to open this encryption or destruction of the information transmitted. The hide of information means include information in the information apparently does not call into doubt and draw attention, and not be aware of before hackers and attackers, so the information will not be rumored to users of the network, but remains content exclusive to the relevant authorities, which are familiar with how to extract this content.

2 Shifting and frequency in encryption

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The ciphertext generated by using keys, the are two types of keys, The first called a similar switch, which uses the same key in the encryption and decryption [1], an example of this type of keys (Caesar) encryption which uses Shifting or the socalled "shift chiper" where each character in the message replaced with the another letter in the alphabet [2], In this type will be very easy to break the key because its used constant key and the character shift in the letter to the another character. For example, when using the 4 key and when there is a character in the letter A to the letter will Encrypts E, its easy to decryption by using the same key [1], This case is easy to break through the use of key characters frequencies by calculate these frequencies in the encrypted message the knowledge of the frequency highest in the ciphertext and returned what is being compensated by the letter E in the original text of the letters that the Supreme frequency and the proportion of its frequency in the texts approximately 12.7, and the character that followed in frequency is the letters t,...Etc, the figure below shows the frequencies of the letters in the language and knowledge of the ciphertext frequencies we can knowledge of the key and break the encryption [3].



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The second type of keys is a different key, used different key encryption from the decryption, an example of this type of algorithm key is (RSA) encryption, which uses two keys a public key and a private key, in this kind scientists were able to get rid of the frequencies of letters and break the encryption and this type of keys to be the toughest in the decoding of the second type, but this kind exhibition vandalized easily by change the public key [2].

3 CLOCK ENCRYPTION

3.1 Encryption Method

Before starting the encryption we will re-arrange all the 96 characters used.

The characters in the standard order of ASCII code is shown in the table below:

 TABLE 1

 ASCII PRINTABLE CHARACTERS+ ENTER / CARRIAGE RETURN CHAR-ACTER [5]

| index | Dec | Hex | Binary | Char- acter | Description |
|-------|-----|-----|----------|----------------|-------------------|
| 0 | 32 | 20 | 00100000 | Space | space |
| 1 | 33 | 21 | 00100001 | ! | exclamation mark |
| 2 | 34 | 22 | 00100010 | 11 | double quote |
| 3 | 35 | 23 | 00100011 | # | number |
| 4 | 36 | 24 | 00100100 | \$ | dollar |
| 5 | 37 | 25 | 00100101 | % | percent |
| 6 | 38 | 26 | 00100110 | & | ampersand |
| 7 | 39 | 27 | 00100111 | 1 | single quote |
| 8 | 40 | 28 | 00101000 | (| left parenthesis |
| 9 | 41 | 29 | 00101001 |) | right parenthesis |
| 10 | 42 | 2A | 00101010 | * | asterisk |
| 11 | 43 | 2B | 00101011 | + | plus |
| 12 | 44 | 2C | 00101100 | , | comma |
| 13 | 45 | 2D | 00101101 | - | minus |
| 14 | 46 | 2E | 00101110 | • | period |
| 15 | 47 | 2F | 00101111 | / | slash |
| 16 | 48 | 30 | 00110000 | 0 | zero |
| 17 | 49 | 31 | 00110001 | 1 | one |
| 18 | 50 | 32 | 00110010 | 2 | two |

| index | Dec | Hex | Binary | Char- acter | Description |
|-------|-----|-----|----------|----------------|---------------|
| 19 | 51 | 33 | 00110011 | 3 | three |
| 20 | 52 | 34 | 00110100 | 4 | four |
| 21 | 53 | 35 | 00110101 | 5 | five |
| 22 | 54 | 36 | 00110110 | 6 | six |
| 23 | 55 | 37 | 00110111 | 7 | seven |
| 24 | 56 | 38 | 00111000 | 8 | eight |
| 25 | 57 | 39 | 00111001 | 9 | nine |
| 26 | 58 | 3A | 00111010 | : | colon |
| 27 | 59 | 3B | 00111011 | ; | semicolon |
| 28 | 60 | 3C | 00111100 | < | less than |
| 29 | 61 | 3D | 00111101 | = | equality sign |
| 30 | 62 | 3E | 00111110 | > | greater than |
| 31 | 63 | 3F | 00111111 | ? | question mark |
| 32 | 64 | 40 | 01000000 | @ | at sign |
| 33 | 65 | 41 | 01000001 | А | |
| 34 | 66 | 42 | 01000010 | В | |
| 35 | 67 | 43 | 01000011 | с | |
| 36 | 68 | 44 | 01000100 | D | |
| 37 | 69 | 45 | 01000101 | Е | |
| 38 | 70 | 46 | 01000110 | F | |
| 39 | 71 | 47 | 01000111 | G | |
| 40 | 72 | 48 | 01001000 | Н | |
| 41 | 73 | 49 | 01001001 | Ι | |
| 42 | 74 | 4A | 01001010 | J | |
| 43 | 75 | 4B | 01001011 | К | |
| 44 | 76 | 4C | 01001100 | L | |
| 45 | 77 | 4D | 01001101 | М | |
| 46 | 78 | 4E | 01001110 | Ν | |
| 47 | 79 | 4F | 01001111 | 0 | |
| 48 | 80 | 50 | 01010000 | Р | |
| 49 | 81 | 51 | 01010001 | Q | |

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| index | Dec | Hex | Binary | Char- acter | Description |
|-------|-----|-----|----------|----------------|-------------------------|
| 50 | 82 | 52 | 01010010 | R | |
| 51 | 83 | 53 | 01010011 | S | |
| 52 | 84 | 54 | 01010100 | Т | |
| 53 | 85 | 55 | 01010101 | U | |
| 54 | 86 | 56 | 01010110 | V | |
| 55 | 87 | 57 | 01010111 | W | |
| 56 | 88 | 58 | 01011000 | Х | |
| 57 | 89 | 59 | 01011001 | Y | |
| 58 | 90 | 5A | 01011010 | Z | |
| 59 | 91 | 5B | 01011011 | [| left square bracket |
| 60 | 92 | 5C | 01011100 | ١ | backslash |
| 61 | 93 | 5D | 01011101 |] | right square bracket |
| 62 | 94 | 5E | 01011110 | ^ | caret / circumflex |
| 63 | 95 | 5F | 01011111 | _ | underscore |
| 64 | 96 | 60 | 01100000 | | grave / accent |
| 65 | 97 | 61 | 01100001 | a | |
| 66 | 98 | 62 | 01100010 | b | |
| 67 | 99 | 63 | 01100011 | С | |
| 68 | 100 | 64 | 01100100 | d | |
| 69 | 101 | 65 | 01100101 | Е | |
| 70 | 102 | 66 | 01100110 | F | |
| 71 | 103 | 67 | 01100111 | g | |
| 72 | 104 | 68 | 01101000 | h | |
| 73 | 105 | 69 | 01101001 | Ι | |
| 74 | 106 | 6A | 01101010 | J | |
| 75 | 107 | 6B | 01101011 | К | |
| 76 | 108 | 6C | 01101100 | L | |
| 77 | 109 | 6D | 01101101 | m | |
| 78 | 110 | 6E | 01101110 | n | |
| 79 | 111 | 6F | 01101111 | 0 | |
| 80 | 112 | 70 | 01110000 | р | |

| index | Dec | Hex | Binary | Char- acter | Description |
|-------|-----|-----|----------|----------------|----------------------------|
| 81 | 113 | 71 | 01110001 | q | |
| 82 | 114 | 72 | 01110010 | R | |
| 83 | 115 | 73 | 01110011 | S | |
| 84 | 116 | 74 | 01110100 | Т | |
| 85 | 117 | 75 | 01110101 | u | |
| 86 | 118 | 76 | 01110110 | V | |
| 87 | 119 | 77 | 01110111 | w | |
| 88 | 120 | 78 | 01111000 | Х | |
| 89 | 121 | 79 | 01111001 | Y | |
| 90 | 122 | 7A | 01111010 | Z | |
| 91 | 123 | 7B | 01111011 | { | left curly bracket |
| 92 | 124 | 7C | 01111100 | | vertical bar |
| 93 | 125 | 7D | 01111101 | } | right curly brack- et |
| 94 | 126 | 7E | 01111110 | ~ | tilde |
| 95 | 13 | 0D | 00001101 | CR | enter / carriage return |

The number of characters that will be used in encryption is 16 characters because we will use the hexadecimal system, for this it must be rearranged with the remaining 80 characters to avoid the juxtaposition between the characters and the purpose of using the largest number of characters when using the shifting indexes of characters as much as possible.

The mechanism of re-arrange characters will be as follows:

- Finding prime numbers between 1 to 96, which are:2,3,5,

- Divide the prime numbers that have been created in the previous point to other groups where we will take the first set of numbers The lowest 10 and we take the second set numbers between 10 and 20 ... etc. as shown in the table below:

- Give the numbers that will be encrypted (1 ... 9) the first number of each group that have been created in the previous point (located in the left hand in each group).

- Give the numbers that will be encrypted (A ... F) the first number of each group that have been created in the second point (from the right) on condition that is the difference between the number and another number at least ten.

The re-arrange the characters will be as in the following table:

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 TABLE 2

 ASCII PRINTABLE CHARACTERS+ ENTER / CARRIAGE RETURN CHARACTER AFTER RE-ARRANGING

| in- dex | Dec | Hex | Binary | Char- acter | Description |
|------------|-----|------------|----------|----------------|-------------------|
| 0 | 32 | 20 | 00100000 | Spac e | space |
| 1 | 33 | 21 | 00100001 | ! | exclamation mark |
| 2 | 48 | 30 | 00110000 | 0 | zero |
| 3 | 34 | 22 | 00100010 | | double quote |
| 4 | 35 | 23 | 00100011 | # | number |
| 5 | 36 | 24 | 00100100 | \$ | dollar |
| 6 | 37 | 25 | 00100101 | % | percent |
| 7 | 70 | 46 | 01000110 | F | |
| 8 | 38 | 26 | 00100110 | & | ampersand |
| 9 | 39 | 27 | 00100111 | 1 | single quote |
| 10 | 40 | 28 | 00101000 | (| left parenthesis |
| 11 | 49 | 31 | 00110001 | 1 | one |
| 12 | 41 | 29 | 00101001 |) | right parenthesis |
| 13 | 42 | 2A | 00101010 | * | asterisk |
| 14 | 43 | 2B | 00101011 | + | plus |
| 15 | 44 | 2C | 00101100 | , | comma |
| 16 | 45 | 2D | 00101101 | - | minus |
| 17 | 46 | 2 E | 00101110 | | period |
| 18 | 47 | 2F | 00101111 | / | slash |
| 19 | 69 | 45 | 01000101 | Е | |
| 20 | 58 | 3A | 00111010 | : | colon |
| 21 | 59 | 3B | 00111011 | ; | semicolon |
| 22 | 60 | 3C | 00111100 | < | less than |
| 23 | 50 | 32 | 00110010 | 2 | two |
| 24 | 61 | 3D | 00111101 | = | equality sign |
| 25 | 62 | 3E | 00111110 | > | greater than |
| 26 | 63 | 3F | 00111111 | ? | question mark |
| 27 | 64 | 40 | 01000000 | @ | at sign |
| 28 | 71 | 47 | 01000111 | G | |

| in- dex | Dec | Hex | Binary | Char- acter | Description |
|------------|-----|-----|----------|----------------|----------------------|
| 29 | 72 | 48 | 01001000 | Н | |
| 30 | 73 | 49 | 01001001 | Ι | |
| 31 | 51 | 33 | 00110011 | 3 | three |
| 32 | 74 | 4A | 01001010 | J | |
| 33 | 75 | 4B | 01001011 | Κ | |
| 34 | 76 | 4C | 01001100 | L | |
| 35 | 77 | 4D | 01001101 | М | |
| 36 | 78 | 4E | 01001110 | Ν | |
| 37 | 68 | 44 | 01000100 | D | |
| 38 | 79 | 4F | 01001111 | 0 | |
| 39 | 80 | 50 | 01010000 | Р | |
| 40 | 81 | 51 | 01010001 | Q | |
| 41 | 82 | 52 | 01010010 | R | |
| 42 | 83 | 53 | 01010011 | S | |
| 43 | 52 | 34 | 00110100 | 4 | four |
| 44 | 84 | 54 | 01010100 | Т | |
| 45 | 85 | 55 | 01010101 | U | |
| 46 | 86 | 56 | 01010110 | V | |
| 47 | 87 | 57 | 01010111 | W | |
| 48 | 88 | 58 | 01011000 | Х | |
| 49 | 89 | 59 | 01011001 | Y | |
| 50 | 90 | 5A | 01011010 | Z | |
| 51 | 91 | 5B | 01011011 | [| left square bracket |
| 52 | 92 | 5C | 01011100 | \ | backslash |
| 53 | 53 | 35 | 00110101 | 5 | five |
| 54 | 93 | 5D | 01011101 |] | right square bracket |
| 55 | 94 | 5E | 01011110 | ^ | caret / circumflex |
| 56 | 95 | 5F | 01011111 | _ | underscore |
| 57 | 96 | 60 | 01100000 | ` | grave / accent |
| 58 | 97 | 61 | 01100001 | a | |
| 59 | 67 | 43 | 01000011 | С | |

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| in- dex | Dec | Hex | Binary | Char- acter | Description |
|------------|-----|-----|----------|----------------|--------------------|
| 60 | 98 | 62 | 01100010 | b | |
| 61 | 99 | 63 | 01100011 | С | |
| 62 | 100 | 64 | 01100100 | d | |
| 63 | 101 | 65 | 01100101 | Е | |
| 64 | 102 | 66 | 01100110 | F | |
| 65 | 103 | 67 | 01100111 | g | |
| 66 | 104 | 68 | 01101000 | h | |
| 67 | 54 | 36 | 00110110 | 6 | six |
| 68 | 105 | 69 | 01101001 | Ι | |
| 69 | 106 | 6A | 01101010 | J | |
| 70 | 107 | 6B | 01101011 | K | |
| 71 | 55 | 37 | 00110111 | 7 | seven |
| 72 | 108 | 6C | 01101100 | L | |
| 73 | 109 | 6D | 01101101 | m | |
| 74 | 110 | 6E | 01101110 | n | |
| 75 | 111 | 6F | 01101111 | 0 | |
| 76 | 112 | 70 | 01110000 | p | |
| 77 | 113 | 71 | 01110001 | q | |
| 78 | 114 | 72 | 01110010 | R | |
| 79 | 66 | 42 | 01000010 | В | |
| 80 | 115 | 73 | 01110011 | S | |
| 81 | 116 | 74 | 01110100 | Т | |
| 82 | 117 | 75 | 01110101 | u | |
| 83 | 56 | 38 | 00111000 | 8 | eight |
| 84 | 118 | 76 | 01110110 | V | |
| 85 | 119 | 77 | 01110111 | w | |
| 86 | 120 | 78 | 01111000 | Х | |
| 87 | 121 | 79 | 01111001 | Y | |
| 88 | 122 | 7A | 01111010 | Z | |
| 89 | 65 | 41 | 01000001 | А | |
| 90 | 123 | 7B | 01111011 | { | left curly bracket |

| in- dex | Dec | Hex | Binary | Char- acter | Description |
|------------|-----|-----|----------|----------------|------------------------------|
| 91 | 57 | 39 | 00111001 | 9 | nine |
| 92 | 124 | 7C | 01111100 | I | vertical bar |
| 93 | 125 | 7D | 01111101 | } | right curly bracket |
| 94 | 126 | 7E | 01111110 | ~ | tilde |
| 95 | 13 | 0D | 00001101 | CR | enter / carriage re- turn |

After rearranging the characters being starting the encryption, which can be summarized as follows:

1. Taking the original text and convert it to the ASCII Code.

2. Convert all the 0 to 1 and vice versa.

3. Took all four bits and converted to hexadecimal system.

4. Find the number of result characters.

5. Taking the first four places of the number of result characters, if the number of resulting characters less than four ranks being as the following:

- If the number consists of three ranks, the one remaining rank compensates 0 (in place of the fourth rank).

- If the number consists of two ranks, the two remaining ranks compensate 0 (in place of the third and fourth ranks).

- If the number consists of one rank, the three remaining ranks compensate 0 (in place of the second, third and fourth ranks).

- If the first and second ranks of the number of result characters more than 12 characters being the result of finding the rest of the division by 12.

- If the third and fourth ranks of the number of result characters more than 60 being the result of finding the rest of the division by 60.

6-The resulting number (from 4,5) converted to the time format (The first and second ranks from left refer to the hours, and the third and fourth ranks from left refer to the minutes).

7- being start encryption by using a clock where Put the time that was obtained in the previous step (point 6), where will indicate the minutes to the number of shifting positions and will start from the first letter, With each character are added minutes, the number of shifting positions increases one by the value of the hour for every adding minute while the times of shifting positions remain fixed until the start of a new hour, as it is shown in the following figure:



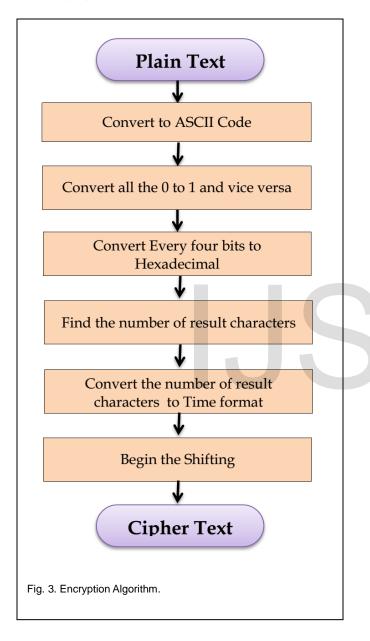
Fig. 2. Example of Cock Encryption.

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8-Obtaining the cipher text.

3.2 Encryption Algorithm

The following figure shows the encryption algorithm after rearranging the characters:



3.3 Example

The Plain Text: Hello

- Convert to ASCII Code: 01001000 01100101 01101100 01101101
- Convert to Hexadecimal: B79A939390
- Convert to Time: 10:00
- Begin the Shifting from 10:01 as

B= index of B+10*1 7=index of 7+10*2 0=index of 9+10*3 A=index of A+10*4 9=index of 9+10*5 3=index of 3+10*6 9=index of 9+10*7 3=index of 3+10*8 9=index of 9+10*9 0=index of 0+10*10 The Cipher Text: A9JKU9g,w%

4 CONCLUSION

Although a huge development in technology, but the encryption in this manner is characterized by the high security addition to the possibility be used manually, it reduces the difference frequency between the characters In addition it is composed of more than one method nested encryption and thus it gives the algorithm for high-security feature and the difficulty of breaking it.

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